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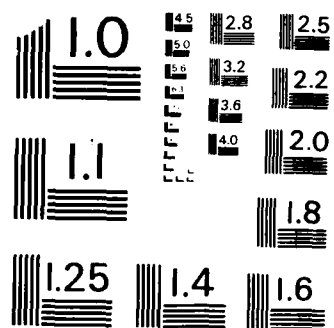
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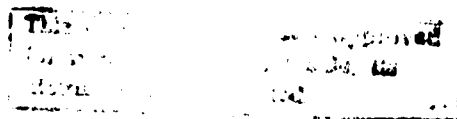
M/G/1 Subject to an Initial
Quorum of Customers

by

Martin Krakowski

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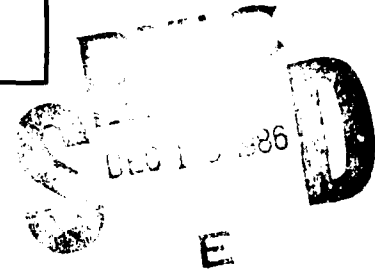
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Abstract

The model $M/G/1$ is modified in the following manner. The server is idled when he runs out of customers and resumes serving when the $(m+1)$ st customer arrives, where $m \geq 0$; $m+1$ is referred to as the "initial quorum" and m is the maximal queue size (and system size) while the server idles or vacations. The modified model is designated as $M/G/1(m)$. When $m=0$ we have the regular $M/G/1$.

In Section 1 we derive the omni-equations for the backlog process B (equation (1.17)) and for the delay process w (equation (1.24)) by exploiting the simple relation between B and w in $M/G/1(m)$, a relation which results in equation (1.12).

In Section 2 we derive several composition properties for the backlog and for the delay. In particular, equation (2.8) says that the backlog is distributed like the sum of the backlog in $M/G/1$ and of a random variable which depends on the service time; and equation (2.9) says that the delay is distributed like the sum of the delay in $M/G/1$ and of a random variable which depends on the service time and on the interarrival time.

Notation

A = interarrival interval

λ = arrival rate = $1/\bar{A}$

x = service duration

μ = $1/\bar{x}$

ρ = λ/μ

r = residual time ("residue") of x ; x and r are related through the omni-equation

$$E\phi(x) - \phi(0) = ExE\phi'(r); \text{ cf. Krakowski (Sept. 1984)}$$

w = delay encountered by a true or virtual customer; in regular $M/G/1$ we have $w=B$

B = backlog (unfinished work) due to incumbent customers; in regular $M/G/1$ we have

$$w = B$$

w_* = delay when server works

B_* = backlog when server works

m = maximal size of the queue when the server idles; $m+1$ is the "initial quorum"

N = size of the system

n = size of the queue

P_j = $\Pr(N=j)$

$p_j = \Pr(n=j)$

$Q_j = \Pr(n=j \text{ and server idles}) , 0 \leq j \leq m ; Q_0 = P_0$

$P_* = \Pr(\text{server works})$

X, Y, and Z are defined by (1.5), (1.6) and (1.23)

$\phi ()$ is an arbitrary well-behaved function of its argument(s); cf. Krakowski (1984, 1985); polynomials and their limits are well-behaved if their expectation is finite; so is the step function; $E\phi(z)$, the expectation of $\phi(z)$, is called the omni-transform of z .

The *omni-convention* calls for taking the expectation of each-side of an omni-equation without explicitly showing the expectation operator; thus $\phi(y) = \phi(z)$ stands for $E\phi(y) = E\phi(z)$

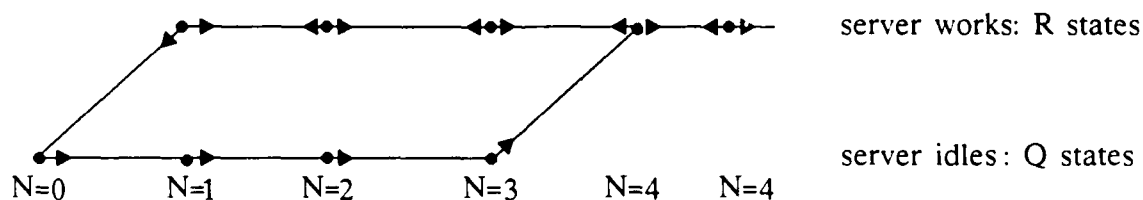
A *free copy* of a random variable z is a random variable distributed like z but independent of z and of any other variable within the scope of the same expectation operator (the expectation may be implied by the omni-convention)

$j*z \stackrel{d}{=} z_1 ++ z_j$ where j is a positive integer and the z_1 are free copies of z (the "generic" random variable); $0*z = 0$

The regular model M/G/1 is modified as follows. The server is idled (or goes off on vacation) each time he runs out of customers, and resumes serving at the instant of the arrival of the $(m+1)$ -th customer, where $m \geq 0$. Thus $m+1$ constitutes an *initial quorum* and m is the maximal queue size while the server idles; we can think of m as a “limbo.” We will designate this model as M/G/1/(m). Thus, when $m = 0$ we have the regular M/G/1.

$$Q_i = (1-\rho)/(m+1) \text{ and } P_s = \rho = \lambda/\mu \quad (1.1)$$
$$\phi(n) = \frac{1-\rho}{m+1} [\phi(0) + \phi(m)] + \rho \phi(n+h) \quad (1.2)$$

Consider the transition diagram for M/G/1/(3):



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In the model $M/G/1/(m)$ the treatment of the delay, the theme of the current report, is more complex than the treatment of the queue-size. It is also more complex and more subtle than the analysis of the delay for the regular $M/G/1$. The instinctive approach in analyzing the delay, in $M/G/1/(m)$ as well as in $M/G/1$, is to establish a balance equation for the virtual delay. This has been a simple task in $M/G/1$, mainly, in our opinion, because in this model the process "virtual delay" and the process "virtual backlog" are identical. In $M/G/1/(m)$ when $m \geq 1$ the two processes differ. It turns out that the analysis of $M/G/1/(m)$ is much simplified, and more insightful, if one analyzes the backlog process B along with the delay process w ; the backlog B is the aggregate remaining service time for all incumbent customers. (Still, the direct "instinctive" approach of balancing the virtual delay, or rather, in the omni-fashion, an arbitrary function of this delay, is a viable approach.) And, of course, the process B may be of interest in its own right and not as a mere stepping stone to analyzing w .

Note The actual delay experienced by a true customer has the same distribution as the virtual delay; cf. Wolff (1982). This follows from the fact that a true customer and a random observer see, stochastically speaking, the same picture when the arrival source is poissonian. Using the same symbol w for true and virtual delays should cause no ambiguity in the context of this report.

The Static Balance A "random" observer will find, with probability Q_0 , that $B = 0$ and $N = 0$; he will find, with probability Q_1 , that $N = 1$ and the server idles and $B = x$; and more generally he will find, with probability Q_j where $0 \leq j \leq m$, that $N = j$ and the server idles $B = j \cdot x$; and he will find with probability P_* that $B = B_*$. Thus, since each $Q_j = (1-\rho)/(m+1)$ and $P_* = \rho$, we have

$$\phi(B) = \frac{1-\rho}{m+1} [\phi(0 \cdot x) + \phi(x) + \phi(2 \cdot x) + \dots + \phi(m \cdot x)] + \rho \phi(B_*) \quad (1.3)$$

When, e.g., $m = 3$ equation (1.3) becomes

$$\phi(B) = \frac{1-\rho}{4} [\phi(0) + \phi(x) + \phi(2 \cdot x) + \phi(3 \cdot x)] + \rho \phi(B_*) \quad (1.3a)$$

To find the omni-equations for the delay w which correspond to (1.3) and (1.3a) we note that a virtual customer who finds, with probability Q_j , j (true) customers in the limbo must await $m - j$ additional customers to reach the quorum of $m + 1$ (this quorum includes the virtual customer) and then must wait until the j senior customers are served. The entire delay of our virtual customer is therefore $(m-j) \cdot A + j \cdot x$, where

$0 \leq j \leq m$; thus the delay of a customer who finds the server idle is composed of m intervals. (When $m = 0$ there is no delay when the server idles, this being the regular M/G/1.) Hence,

$$\phi(w) = \frac{1-\rho}{m+1} [\phi(m \cdot A) + \phi((m-1) \cdot A + x) + \dots + \phi(m \cdot x)] + \rho \phi(w_*) \quad (1.4)$$

When $m = 3$ equation (1.4) becomes

$$\begin{aligned} \phi(w) = \frac{1-\rho}{4} [\phi(3 \cdot A) + \phi(2 \cdot A + x) + \phi(A + 2 \cdot x) + \phi(3 \cdot x)] \\ + \rho \phi(w_*) \end{aligned} \quad (1.4a)$$

where w_* is the delay conditioned upon the server working.

We simplify the typography of (1.4) and (1.4a) by defining

$$\phi(X) \triangleq \frac{1}{m+1} \sum_{j=0}^m \phi(j \cdot x) \quad (1.5)$$

and

$$\phi(Y) \triangleq \frac{1}{m+1} \sum_{j=0}^m \phi((m-j) \cdot A + j \cdot x) \quad (1.6)$$

Then (1.3) and (1.4) become

$$\phi(B) = (1-\rho) \phi(X) + \rho \phi(B_*) \quad (1.7)$$

and

$$\phi(w) = (1-\rho) \phi(Y) + \rho \phi(w_*) \quad (1.8)$$

A key observation in our approach is that in the model M/G/1/(m) the virtual delay and the backlog observed at an instant when the server works are one and the same process; hence

$$B_* = w_* \quad (1.9)$$

from which follows the weaker statement that

$$\phi(B_*) = \phi(w_*) \quad (1.10)$$

which merely says that B_* and w_* are identically distributed. Hence, from (1.8) and (1.10) it follows that

$$\phi(w) = (1-\rho) \phi(Y) + \rho \phi(B_*) \quad (1.11)$$

It follows from (1.7) and (1.11) that

$$\phi(w) = \phi(B) + (1-\rho) [\phi(Y) - \phi(X)] \quad (1.12)$$

Thus if we have an omni-equation for B we can transpose it with the aid of (1.12) into an omni-equation for w . We now derive such an equation as balancing condition for an arbitrary function of B .

The backlog B jumps to $B + x$ with each arrival, these arrivals being of frequency λ . (Departures of serviced customers do not affect the balance since in our model we assume that the value of a customer's service time is revealed at the instant of his arrival.) On the other hand, while the server works, with probability P_* , he keeps on reducing the backlog now denoted by B_* at the rate $dB_* = dt$. Hence the omni-balance for B is

$$\lambda [\phi(B+x) - \phi(B)] = P_* \phi'(B_*) \quad (1.13)$$

According to the shifted renewal equation (Krakowski, September 1984)

$$\phi(B+x) - \phi(B) = \bar{x} \phi'(B+r) \quad (1.14)$$

where r is the residue of x ; the differentiation is with respect to the entire argument. From (1.13) and (1.14) we have

$$\lambda \bar{x} \phi'(B+r) = P_* \phi'(B_*)$$

which integrated typographically is

$$\rho \phi(B+r) = P_* \phi(B_*) \quad (1.15)$$

Putting $\phi(\cdot) = 1$ we rederive $P_* = \rho$ and (1.15) becomes

$$\phi(B_*) = \phi(B+r) \quad (1.16)$$

From (1.17) and (1.16) we have

$$\phi(B) = (1-\rho) \phi(X) + \rho \phi(B+r) \quad (1.17)$$

the sought for omni-equation for the backlog B . From (1.17) we can get, recursively, successive moments of B and a convolution equation for the distribution of B . And in turn, with the aid of (1.12), we can get the successive moments of w and a convolution equation for the distribution of w . Of course, all needed moments and distributions of X and Y and Z have to be derived as a side exercise.

Example: Find \bar{B} and \bar{w}

From (1.15) we have

$$\bar{X} = \frac{1}{m+1} [\bar{x} + m\bar{x}] = \frac{m(m+1)\bar{x}}{2(m+1)} = \frac{1}{2} m\bar{x} \quad (1.18)$$

and from (1.6) we have, similarly,

$$\bar{Y} = \frac{1}{2} m\bar{A} + \frac{1}{2} m\bar{x} = \frac{1}{2} m\left(\frac{1}{\lambda} + \frac{1}{\mu}\right) \quad (1.19)$$

From (1.17) and (1.18) we have

$$\bar{B} = \frac{1}{2} m\bar{x} + \frac{\rho\bar{r}}{1-\rho} \quad (1.20)$$

From (1.12), (1.18) and (1.19) we have

$$\begin{aligned} \bar{w} &= \bar{B} + (1-\rho)(\bar{Y} - \bar{X}) = \frac{1}{2} m\bar{x} + \frac{\rho\bar{r}}{1-\rho} + (1-\rho)m\bar{A} \\ &= \frac{m}{2\lambda} + \frac{\rho\bar{r}}{1-\rho} \end{aligned} \quad (1.21)$$

Thus, \bar{w} is composed of m/λ and the expected waiting time for a regular M/G/1. This is not incidental and we will take up the problem of composition properties in the model M/G/1(m) in Section 2.

To complete the current section we derive the omni-equation for w . From (1.12) and (1.17) we have

$$\phi(w) = (1-\rho)\{\phi(Y) - \rho[\phi(Y+r) - \phi(X+r)]\} + \rho\phi(w+r) \quad (1.22)$$

A comparison of equation (1.22) with the quite simpler (1.17)—even if we replace $\phi(X)$ by its explicit representation in (1.5)—will give the reader an idea about the more extensive need for algebraic maneuvering in a direct derivation of (1.22) from the omni-balance of the virtual delay.

The formal appearance of (1.21) is simplified by defining

$$\phi(Z) \triangleq \phi(Y) - \rho[\phi(Y+r) - \phi(X+r)] \quad (1.23)$$

From (1.22) and (1.23) we have

$$\phi(w) = (1-\rho)\phi(Z) + \rho\phi(w+r) \quad (1.24)$$

It is easy to verify that (1.24) leads to \bar{w} as given in (1.21).

Section 2: Composition Theorems

When $m = 0$, i.e. when we consider the regular $M/G/1/(0)$ alias $M/G/1$, we have according to (1.5), (1.6) and (1.23)

$$\phi(X_0) = 0, \quad \phi(Y_0) = 0, \quad \phi(Z_0) = 0, \quad (2.1)$$

We will use in this section, as we have just done in (2.1), a subscript serving as a reminder of the size of m in the model. This is not a necessity, just a convenience; most equations in this section relate a general $M/G/1/(m)$ model to the regular $M/G/1/(0)$ so that the reminder-subscripts will identify the models. Thus we will write X_m , Y_m , Z_m , B_m , and w_m , in place of X , Y , Z , B and w . Equations (1.17) and (1.22) for $m = 0$ are written now, taking into account (2.1),

$$\phi(B_0) = (1-\rho)\phi(0) + \rho\phi(B_0+r) \quad (2.2)$$

and

$$\phi(w_0) = (1-\rho)\phi(0) + \rho\phi(w_0+r) \quad (2.3)$$

Equations (2.2) and (2.3) imply, as can be shown, that

$$\phi(B_0) = \phi(w_0) \quad (2.4)$$

which we know from the structure of the regular $M/G/1$.

Let us now shift (2.2) by X_m thus obtaining

$$\phi(B_0 + X_m) = (1-\rho)\phi(X_m) + \rho\phi(B_0+r+X_m) \quad (2.5)$$

From (2.5) and (1.17) we have

$$\phi(B_m) - \rho\phi(B_m+r) = \phi(B_0+X_m) - \rho\phi(B_0+r+X_m) \quad (2.6)$$

Defining

$$\Psi(B_m) \triangleq \phi(B_m) - \rho\phi(B_m+r) \quad (2.7)$$

we write (2.6) as

$$\Psi(B_m) = \Psi(B_0 + X_m)$$

or, reverting to $\phi(\cdot)$ in place of $\Psi(\cdot)$, as

$$\phi(B_m) = \phi(B_0 + X_m) \quad (2.8)$$

Equation (2.8) states that the backlog B_m in $M/G/1/(m)$ is distributed like the convolution of B_0 , the backlog in $M/G/1$, and X_m .

From (2.3) and (1.24) we derive in like manner

$$\phi(w_m) = \phi(w_0 + Z_m) \quad (2.9)$$

thus showing that the delay in $M/G/1/(m)$ is distributed like the convolution of w_0 , the delay in regular $M/G/1$, and Z_m .

The composition equations (2.8) and (2.9) should be especially useful computationally when the distribution or moments of the delay w_0 are known and when m is treated as a parameter. Of course, the distribution, or moments, of Z_m have to be evaluated as a side exercise. Similar remarks hold for the distribution and moments of B should the backlog be a process of interest. The labor involved may be formidable but in some practical applications it might be justified economically.

Another composition property for the delay w_m is obtained as follows. Providing (1.12) with m -subscripts we have

$$\phi(w_m) = (1-\rho)[\phi(Y_m) - \phi(X_m)] + \phi(B_m) \quad (2.10)$$

Since in the regular $M/G/1$, i.e. in $M/G/1/(0)$, $B_0 = w_0$ we can write (2.8) in the form

$$\phi(B_m) = \phi(w_0 + X_m) \quad (2.11)$$

From (2.10) and (2.11) we have

$$\phi(w_m) = (1-\rho)[\phi(Y_m) - \phi(X_m)] + \phi(w_0 + X_m) \quad (2.12)$$

Equation (2.12) appears simpler numerically than (2.9), especially for the computation of moments.

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